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# Set colorings of graphs

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#### ABSTRACT

A **set coloring** of the graph *G* is an assignment (function) of distinct subsets of a finite set *X* of **colors** to the vertices of the graph, where the colors of the edges are obtained as the symmetric differences of the sets assigned to their end vertices which are also distinct. A set coloring is called a **strong set coloring** if sets on the vertices and edges are distinct and together form the set of all nonempty subsets of *X* are obtained on the edges. A graph is called **strongly set colorable** (properly set colorable) if it admits a strong set coloring (proper set coloring).

In this paper we give some necessary conditions for a graph to admit a strong set coloring (a proper set coloring), characterize strongly set colorable complete bipartite graphs and strongly (properly) set colorable complete graphs, etc. Also, we give a construction of a planar strongly set colorable graph from a planar graph, a strongly set colorable tree from a tree and a properly set colorable tree from a tree, etc., thereby showing their embeddings. © 2008 Elsevier Ltd. All rights reserved.

#### 1. Introduction

In this paper we consider only finite simple graphs. For all notation in graph theory we follow Harary [3] and West [5].

Colorings of the vertices and edges of a graph *G* which are required to obey certain conditions have often been motivated by their utility in various applied fields and their intrinsic mathematical interest (logico-mathematical). An enormous amount of literature has built up on several kinds of colorings of graphs.

Motivated by the papers of Hopkroft and Krishnamurty [4], Balister et al. [2], and Acharya [1], we introduce set colorings of graphs: Let X be a nonempty set of colors,  $2^X$  denote the set of all possible

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Fig. 1. Set coloring of a Heawood graph.

combinations of colors (or power set) of *X* and  $Y(X) = 2^X \setminus \{\emptyset\}$ . For any two subsets *A* and *B* of *X* let  $A \oplus B$  denote the symmetric difference of *A*, *B* and be given by  $A \oplus B = (A \cup B) - (A \cap B)$ .

Given a (p, q)-graph G = (V, E) and a nonempty set X of colors, we define a function f on the vertex set V of G as an assignment of subsets of X to the vertices of G, and given such a function f on the vertex set V we define  $f^{\oplus}$  on the set of edges E as an assignment of the colors  $f^{\oplus}(e) = f(u) \oplus f(v)$  to the edge e = uv of G.

Let  $f(G) = \{f(u) : u \in V\}$  and  $f^{\oplus}(G) = \{f^{\oplus}(e) : e \in E\}.$ 

We call f a **set coloring** of G if both f(G) and  $f^{\oplus}(G)$  are injective functions. A graph is called **set colorable** if it admits a set coloring. A set coloring f of G is called a **strong set coloring** if f(G) and  $f^{\oplus}(G)$  are disjoint subsets of X and, further, they form a partition of Y(X). If G admits such a coloring then G is a called a **strongly set colorable graph**.

A set coloring f is called a **proper set coloring** if  $f^{\oplus}(G) = Y(X)$ . If a graph G admits such a set coloring then it is called a **properly set colorable graph**.

The **set coloring number**  $\sigma(G)$  of a graph *G* is the least cardinality of a set *X* with respect to which *G* has a set coloring. Further, if  $f : V \to 2^X$  is a set coloring of *G* with  $|X| = \sigma(G)$  we call *f* an optimal set coloring of *G*.

Theorem 1. For any graph G,

 $\lceil \log_2(q+1) \rceil \le \sigma \ (G) \le p-1,$ 

where  $\lceil x \rceil$  denotes the least integer not less than the real number x, and the bounds are best possible.

Fig. 1 gives an optimal set coloring of a Heawood graph.

#### 2. Strongly (properly) set colorable graphs

Since all the nonempty subsets have to appear in any strong set coloring of a (p, q)-graph G, a necessary condition for G to be strongly set colorable is that  $p + q + 1 = 2^m$ , for the positive integer m = |X|. This necessary condition immediately yields that no cycle is strongly set colorable. Also, we observe that the above condition is not sufficient for saying that a graph G is strongly set colorable





as the path of length 3 satisfies the condition for m = 3 but one can verify that it is not strongly set colorable. We propose:

**Conjecture 1.** No path of length greater than 2 is strongly set colorable.

Similarly, a necessary condition for *G* to be properly set colorable is that  $q + 1 = 2^m$ . From the above necessary condition it follows that the cycles of lengths not equal to  $2^m - 1$  are not properly set colorable.

Fig. 2 gives examples of properly, strongly, non-strongly and non-properly colorable graphs.

The following result gives a natural link between strongly set colorable and properly set colorable graphs.

**Theorem 2.** A graph G is strongly set colorable if and only if  $G + K_1$  with  $V(K_1) = \{v\}$  has a proper set coloring F such that  $F(v) = \emptyset$ .

**Proof.** Let *f* be a strong set coloring of *G*. Then, extend *f* to the vertices of  $G + K_1$  to a set assignment *F* so that the restriction map F|V(G) of *F* to V(G) is *f* and  $F(v) = \emptyset$ . Since *f* is a strong set coloring of *G* the edges of  $G + K_1$  having the form uv where  $u \in V(G)$  will receive f(u). So *F* turns out to be a required proper set coloring of  $G + K_1$ .  $\Box$ 

Conversely, if  $H = G + K_1$  has a proper set coloring F with  $F(v) = \emptyset$  then the removal of v from H obviously results in a strong set coloring of G.

The following two results give stronger necessary conditions for a strong set coloring of graphs.

**Theorem 3.** If a graph G(p > 2) has:

- (i) exactly one or two vertices of even degree or
- (ii) exactly three vertices of even degree, say,  $v_1$ ,  $v_2$ ,  $v_3$ , and any two of these vertices are adjacent or
- (iii) exactly four vertices of even degree, say,  $v_1$ ,  $v_2$ ,  $v_3$ ,  $v_4$  such that  $v_1v_2$  and  $v_3v_4$  are edges in G, then G is not strongly set colorable.

**Proof.** Let *G* be a graph with a strong set coloring *f* with respect to a set *X* having *m* colors. Let  $\{v_1, v_2, v_3, \ldots, v_p\}$  be the vertices of *G* such that  $f(v_i) = A_i$ ,  $1 \le i \le p$ , and  $A_i \in Y(X)$ . Then we get  $f(G) \cup f^{\oplus}(G) = \{A_1, A_2, \ldots, A_p, \{A_i \oplus A_j : v_i v_j \in E\}\} = Y(X)$ . As the symmetric difference of all the nonempty subsets of any set is the empty set, we get

$$f(G) \cup f^{\oplus}(G) = \{A_1, A_2, \dots, A_p\} \cup \{A_i \oplus A_j : v_i v_i \in E\} = \emptyset.$$

$$\tag{1}$$

One can see that if the degree of a vertex v is even then the set A assigned to v appears an odd number of times and if the degree of a vertex u is odd then the set B assigned to u appears an even number of times in (1).

- (i) Suppose that *G* has exactly one vertex of even degree, say  $v_1$ . Then  $A_1$  will appear an odd number of times and all the other sets will appear an even number of times in (1). So, as the binary operation  $\oplus$  is commutative, when the symmetric differences of the sets in (1) are taken, all the subsets which are assigned to the vertices of odd degree will vanish and hence we get that  $A_1 = \emptyset$ , a contradiction to the definition of the strong set coloring of *G*. Hence, if *G* has exactly one vertex of even degree then *G* is not strongly set colorable. Suppose that *G* has exactly two vertices of even degree, say  $v_1, v_2$ . Then using arguments similar to those above and from (1) we obtain that  $A_I \oplus A_2 = \emptyset$ , which implies that  $A_I = A_2$ , a contradiction to the injectivity of *f*. Hence, if *G* has exactly two vertices of even degree then *G* is not strongly set colorable.
- (ii) Suppose *G* has three vertices of even degree, say,  $v_1$ ,  $v_2$ ,  $v_3$  such that  $v_1v_2$  is an edge in *G*. Then by arguments similar to those for (i) and from (1) we obtain that  $A_3 = A_1 \oplus A_2$ , or  $A_2 \oplus A_3 = A_1$ , or  $A_l \oplus A_3 = A_2$ , a contradiction to the definition of a strong set coloring of *G*. Hence, if *G* has exactly three vertices of even degree as mentioned in Theorem 1, then *G* is not strongly set colorable.
- (iii) Suppose that *G* has exactly four vertices of even degree as given in the statement of Theorem 3. Then, by similar arguments and using (1), we get  $A_1 \oplus A_2 = A_3 \oplus A_4$  or  $A_1 \oplus A_3 = A_2 \oplus A_4$  or  $A_1 \oplus A_4 = A_2 \oplus A_3$  respectively, a contradiction to the injectivity of  $f^{\oplus}$ . Hence, if *G* has exactly four vertices of even degree then *G* is not strongly set colorable.  $\Box$

**Theorem 4.** If a graph *G* has a strong set coloring *f* with respect to a set *X* of cardinality *m*, there exists a partition of the vertex set *V* into two sets  $V_1$  and  $V_2$  such that the number of edges joining the vertices of  $V_1$  with those of  $V_2$  is exactly  $2^{m-1} - |V_2|$ .

**Proof.** Suppose that *G* is strongly set colorable with respect to a set *X* of cardinality  $m \ge 2$ . Consider a partition of *V* into two sets  $V_1$  and  $V_2$  such that  $V_1 = \{u \in V : |f(u)| \text{ is even}\}$  and  $V_2 = \{v \in V : |f(v)| \text{ is odd}\}$ . One can obtain other odd subsets of *X* which are not there on the vertices only by taking the symmetric differences between the vertices of  $V_1$  with those of  $V_2$  and hence the result follows, as there are exactly  $2^{m-l}$  subsets of each parity for a set *X* of cardinality *m*.  $\Box$ 

The following two results give stronger necessary conditions for a proper set coloring of graphs.

#### **Theorem 5.** *If* a graph G(p > 2) has:

- (i.) exactly two vertices of odd degree **or**
- (ii.) exactly four vertices of odd degree, say,  $v_1v_2$ ,  $v_3$ ,  $v_4$  with  $v_1v_2$  and  $v_3v_4$  being edges in G, then G is not properly set colorable.

**Proof.** Let *G* have a proper set coloring *f* with respect to a set *X* of cardinality *m*. Let  $\{v_1, v_2, v_3, ..., v_p\}$  be the vertices of *G* such that  $f(v_i) = A_i$ ,  $1 \le i \le p$ , and  $A_i \in Y(X)$ . As *G* is properly set colorable, we have

$$f^{\oplus}(G) = \{A_i \oplus A_j : v_i v_j \in E\} = Y(X).$$
<sup>(2)</sup>

- (i) Suppose that *G* has exactly two vertices of odd degree, say,  $v_1, v_2$ . From (1) we obtain  $A_l \oplus A_2 = \emptyset$ , which implies that  $A_l = A_2$ , a contradiction to the injectivity of *f*. Hence if *G* has exactly two vertices of odd degree then it is not properly set colorable.
- (ii) Suppose that *G* has exactly four vertices of odd degree as mentioned in the theorem. Then by similar arguments we obtain

 $A_1 \oplus A_2 = A_3 \oplus A_4$  or  $A_1 \oplus A_3 = A_2 \oplus A_4$  or  $A_1 \oplus A_4 = A_3 \oplus A_2$  respectively, a contradiction to the injectivity of  $f^{\oplus}$ . Hence, if *G* has exactly four vertices of odd degree as mentioned in the theorem then *G* is not properly set colorable.  $\Box$ 

**Corollary 5.1.** No path of length greater than 2 is properly set colorable.

**Theorem 6.** If a graph *G* has a proper set coloring *f* with respect to a set *X* of cardinality *m*, then there exists a partition of the vertex set *V* into two sets  $V_1$  and  $V_2$  such that the number of edges joining the vertices of  $V_1$  with those of  $V_2$  is exactly  $2^{m-1}$ .

**Proof.** Suppose that *G* is properly set colorable with respect to a set *X* of cardinality  $m \ge 2$ . Let  $V_1$  and  $V_2$  be a partition of *V* as mentioned in the proof of Theorem 4. As one can obtain all the odd subsets of *X* by taking the symmetric differences between the vertices of  $V_1$  with those of  $V_2$ , the result follows.

The next result characterizes the strongly colorable complete graphs.

Consider the complete graph  $K_n$ . Suppose that it is strongly set colorable with respect to a set X of cardinality m. Then it follows that the sum of the number of vertices and the edges of  $K_n$  must be equal to  $2^m - 1$ , i.e.,  $n + n(n - 1)/2 = 2^m - 1$ , which yields the quadratic equation

$$n^{2} + n - (2^{m+1} - 2) = 0.$$
(3)

Solving (3) for positive integer values of *n*, we get

$$n = (1/2)(\sqrt{2^{m+3} - 7}) - 1).$$
(4)

The first four values of n for which  $K_n$  may possibly be strongly set colorable are 1, 2, 5 and 90, where the values of m are 1, 2, 4 and 12, respectively. The following result characterizes the strongly set colorable complete graphs.

#### **Theorem 7.** The nontrivial complete graph $K_n$ is strongly set colorable if and only if n = 2, 5.

**Proof.** Suppose that *G* is *strongly set colorable* with respect to a set *X* of cardinality  $m \ge 2$ . Consider a partition of *V* into two sets  $V_1$  and  $V_2$  as mentioned in the proof of Theorem 4. Then one can obtain all the other odd subsets of *X* which are not covered by  $V_2$  by taking the symmetric differences between the vertices of  $V_1$  and those of  $V_2$ . Thus,  $|V_1||V_2| = 2^{(m-1)} - |V_2|$ ,

i.e., 
$$(|V_1| + 1) |V_2| = 2^{(m-1)}$$
. (5)

Hence,  $|V_1| + 1$  is a power of 2 and  $|V_2|$  is a power of 2. By the symmetric differences among the vertices of  $V_1$  as well as among the vertices of  $V_2$ , we obtain all the other even nonempty subsets of X which are not covered by  $V_1$ . Thus, we get

$$|V_1|(|V_1| - 1)/2 + |V_2|(|V_2| - 1)/2 = 2^{(m-1)} - |V_1| - 1$$
(6)

i.e., 
$$|V_1|(|V_1|+1)/2 + |V_2|(|V_2|-1)/2 = 2^{(m-1)} - 1.$$
 (7)

Eq. (7) implies that one of the terms in (7) is odd, say  $|V_1|(|V_1|+1)/2$  is odd. We know that  $(|V_1|+1)$  is a power of 2. Suppose that  $(|V_1|+1) = 2^t$ ,  $t \ge 2$ ; then we obtain that  $|V_1|(|V_1|+1)/2$  is even, a contradiction. Hence,  $(|V_1|+1) = 2$ . Thus, we obtain  $|V_1| = 1$  and hence,  $|V_2| = n - 1$ . Then, from (5) we get

$$2(n-1) = 2^{(m-1)}.$$
(8)

Also, from (7), we obtain

$$2 + (n-1)(n-2)/2 = 2^{(m-1)}.$$
(9)

Similarly, if  $|V_2|(|V_2| - 1)/2$  is odd, then (8) and (9) are again obtained. Equating (8) and (9), we obtain

2(n-1) = 2 + (n-1)(n-2)/2 or  $n^2 - 7n + 10 = 0$ , which implies that n = 2, 5.

Conversely, suppose that n = 2, 5; then one can easily verify that  $K_2$  and  $K_5$  are strongly set colorable.  $\Box$ 

**Theorem 8.** The complete graph  $K_n$  is properly set colorable with respect to a set X of cardinality m if and only if n = 2, 3 and 6.

**Proof.** The proof follows from Theorems 7 and 2.  $\Box$ 

**Theorem 9.** The nontrivial complete n-ary tree  $T_n^t$  is strongly set colorable if and only if  $n = 2^{\alpha} - 1$  and t = 1, where t is the number of levels of  $T_n^t$ .

**Proof.** Suppose that  $G = T_n^t$  is strongly set colorable with respect to a set *X* of cardinality *m*. Then we obtain  $|V(G)| + |E(G)| = 2^m - 1$ .

The case when n is even follows from Theorem 3. Thus, no complete n-ary tree G is strongly set colorable when n is even.

By the definition of a complete *n*-ary tree, we obtain

 $(1 + n + n^2 + \dots + n^t) + (1 + n + n^2 + \dots + n^t - 1) = 2^m - 1$  or  $(2n^{(t+1)} - n - 1)/(n - 1) = 2^m - 1$  or  $(n^{(t+1)} - 1)/(n - 1) = 2^{(m-1)}$ ,

which implies that *n* is odd and hence *t* is odd. Thus, from  $(n^{(t+1)} - 1)/(n-1) = 2^{(m-1)}$ , we obtain  $(1 + n + n^2 + \cdots + n^t) = 2^{(m-1)}$  or

$$(1+n)(1+n^2+n^4+\dots+n^{t-1})=2^{(m-1)}$$
(10)

which implies that  $(1 + n) = 2^{\alpha}$ ,  $\alpha$  is a positive integer. Thus, from (10) we obtain

$$(1+n^2+n^4+\dots+n^{t-1})=2^{(m-\alpha-1)}.$$
(11)

One can write (11) as

$$(1+n^2)(1+n^4+n^8+\cdots+n^{t-3})=2^{(m-\alpha-1)}$$

which implies that  $1 + n^2 = 2^{\beta}$ . Substituting the value of *n* from  $(1 + n) = 2^{\alpha}$ , we obtain

 $1 + (2^{\alpha} - 1)^{2} = 2^{\beta}, \text{ or }$   $2^{2\alpha} - 2^{\alpha+1} + 2 = 2^{\beta}, \text{ or }$   $2^{2\alpha-1} - 2^{\alpha} + 1 = 2^{\beta-1}, \text{ or }$  $2^{2\alpha-1} - 2^{\beta-1} = 2^{\alpha} - 1$ 

which implies that  $2^{\alpha} - 1$  is even, or  $\alpha = 0$ , or n = 0, a contradiction. Thus,  $1 + n = 2^{\alpha}$  or  $n = 2^{\alpha} - 1$ . Also, from (11) we obtain  $1 = 2^{(m-\alpha-1)}$  or  $m = \alpha + 1$  and also t = 1.

Conversely, suppose that  $n = 2^{\alpha} - 1$  and t = 1. Then *G* reduces to the star  $K_{1,2^{\alpha}-1}$ . Let  $X = \{1, 2, ..., m\}, X_1 = \{1\}$  and  $X_2 = \{2, 3, ..., m\}$ . Assign the set  $X_1$  to the central vertex and all the nonempty subsets of  $X_2$  to the remaining vertices of the star in a one-to-one manner. Then it is not hard to verify that the assignment is a *strong set coloring* of  $K_{1,2^{\alpha}-1}$ .

A similar proof proves the following theorem.

**Theorem 10.** The nontrivial complete n-ary tree *G* is properly set colorable if and only if  $n = 2^{\alpha} - 1$  and t = 1.

**Theorem 11.** The complete bipartite graph  $K_{a,b}$  is strongly set colorable if and only if  $(a+1)(b+1) = 2^m$ , where *m* is a positive integer.

**Proof.** Let  $K_{a,b}$  be strongly set colorable with respect to a set *X* of cardinality *m*. Then it follows that  $|V(K_{a,b})| + |E(K_{a,b})| = 2^m - 1$ , i.e.,  $a + b + ab = 2^m - 1$ , which yields

 $(a+1)(b+1) = 2^m$ .

Conversely, assume that

$$(a+1)(b+1) = 2^m,$$
(12)

for some positive integers a, b and m, where m is the cardinality of the set X. Taking the logarithm to base 2 on both sides of (12), we obtain

 $m = \log_2(a+1) + \log_2(b+1).$ 

Hence, there exists a partition  $\{X_1, X_2\}$  of X such that  $|X_1| = \log_2(a + 1)$  and  $|X_2| = \log_2(b + 1)$ . Let  $A_1$  and  $A_2$  constitute the bipartition of the vertex set of  $K_{a,b}$ . Assign the nonempty subsets of  $X_i$  to the vertices in  $A_i$ , i = 1, 2, in a one-to-one manner. Then one can verify that the resulting assignment is indeed a strong set coloring of  $K_{a,b}$ .  $\Box$ 

**Conjecture 2.** The complete bipartite graph  $K_{a,b}$  is properly set colorable if and only if it is a star with a = 1 and  $b = 2^{n-1}$ .

Next, we give some results on the construction of strongly (properly) set colored graphs and show their embeddings.

Let G be the given planar graph with n vertices. Let T be a spanning tree of G. Introduce a new vertex v and join it to a vertex of G which is in the exterior face. As T is a spanning tree, let the new tree with v as the additional vertex be the tree  $T_l$ . Draw the tree  $T_l$  as a rooted tree with root v. Let  $X = \{1, 2, \dots, n\}$  be a set of cardinality n. Let the vertices of T be  $v_i$ , which are in ascending order in T<sub>i</sub>. Assign the set X to v and the single-element subsets  $\{i\}$  of X to the remaining vertices  $v_i$  such that  $f(v_i) = \{i\}$ . Let  $\{A_{i,j} : v_i v_j \in E, i < j\}$  be the t two-element subsets of X which are already obtained on the edges of G and  $B_1, B_2, \ldots, B_k$  be the remaining two-element subsets of X such that t + k = (n - 1)(n)/2. We know that (n - 2)-element subsets are the complements of 2-elements, (n-3)-element subsets are the complements of 3-elements, ..., (n-2)/2-element subsets are the complements of (n + 2)/2-elements if n is even and (n - 1)/2-element subsets are the complements of (n + 1)/2-elements if n is odd. Hence introducing the required number of new vertices, joining them to the vertex v and assigning the subsets of cardinality (n-3), (n-4), (n-5), etc., up to (n-1)/2-element subsets, if n is odd (up to half of the (n/2)-element subsets, if n is even), we can obtain all the subsets of cardinality 3, 4, 5, ..., (n-1)/2 if n is odd (and (n/2) if n is even). Thus, we have exhausted all the subsets except (n(n-1)/2 - t), two-element subsets, (n-2)-element subsets and (n-1), (n-1)-element subsets.

Introduce *k* new vertices and join them to the vertex *v*. Then assign the sets  $B_1, B_2, \ldots, B_k$  to these newly introduced vertices in a one-to-one manner. This assignment will generate the remaining *k* two-element subsets on these new edges. Thus, we have covered *X*, all the single-element subsets, one (n - 1)-element subset and all the two-element subsets of *X*. We have to obtain the remaining (n - 1)(n - 1)-element subsets and the remaining (n - 2)-element subsets, either on vertices or edges.

To generate (n - 1)-element subsets, consider the internal vertices of  $T_1$  (internal vertices are the vertices of degree at least 2). Starting from the first layer of T, whenever  $v_i v_j$  is an edge, introduce a new vertex and assign the complement of the set  $A_{i,j}$ , and join the new vertex to the internal vertex of  $T_l$  which was assigned the single-element subset containing the element *i*, which will yield a (n - 1)-element subset on the new edge. Continue the procedure until the last but one layer and until all the internal vertices of  $T_1$  are exhausted. Thus, we have generated the (n - 1)-, (n - 2)-element subsets and (n - 1)-element subsets of X. Still we have to cover (t - (n - 1)), (n - 2)-element subsets. Introduce t + 1 - n new vertices and assign the remaining (t - (n - 1)) (n - 2)-element subsets to these isolated vertices. Thus the resulting graph  $G_l$  is strongly set colorable and planar.



Fig. 3. Illustration of the procedure in the construction.

**Remark 2.** From the construction in Fig. 3 it follows that every planar graph can be embedded as an induced subgraph of a strongly set colored planar one. Also, it follows that any tree can be embedded as an induced subgraph of a strongly set colorable tree, as there will not be any isolated vertices if *G* itself is a tree.

Given below is a construction of a bigger properly set colored tree from a properly set colored tree. Suppose that a tree *T* is properly set colored with respect to a set *X* of cardinality *m*. Then all the  $2^m$  subsets of *X* appear on the vertices of *T*. Introduce  $2^n - 1$  isolated vertices and join them to the vertices of *T*. Assign the nonempty subsets of a set  $X'(X \cap X' = \emptyset)$  of cardinality *n* to the newly introduced vertices in a one-to-one manner. Then it is not hard to verify that all the  $2^{m+n} - 1$  nonempty subsets of the set  $Y = X \cup X'$  of cardinality m + n will appear on the edges of the resulting graph. Hence the resulting graph is properly set colored.

The construction Fig. 4 proves that every tree can be embedded as an induced subgraph of a properly set colored tree.

Let *T* be the tree with *n* vertices. We prove the result by induction on the number of edges of *T*. One can easily see that trees with one or two edges can be embedded as an induced subgraph of a properly set colorable tree. Suppose that the result is true for a tree  $T_1$  with n - 2 edges, where  $T_1$  is obtained from *T* by removing a pendant edge uv such that v is in  $T_1$ . This means a tree with n - l vertices and n - 2 edges can be embedded as an induced subgraph of a properly set colorable tree, say,  $T_2$ .

Let X be the set of cardinality m with respect to which  $T_2$  has a proper set coloring f. Since  $T_2$  is a tree,  $f(T_2) = 2^X$ . Then join the edge uv. Add an element y to all the  $2^X$  sets which are assigned to the vertices of  $T_2$  and assign the set Ø to the vertex u of T. Let  $f(v) = S \subset X$  (where S is a subset of X). Note the set  $S \cup \{y\}$  has been assigned to v. The set  $S \cup \{y\}$  is obtained on the edge uv. Then introduce a new vertex and join it to v. Assign the set S to the newly introduced vertex and then the set  $\{y\}$  is generated on the new edge. Let  $X_I = X \cup \{y\}$ . Introduce  $2^m - 2$  new vertices and join them to the vertex w where  $f(w) = \emptyset$  and assign all the elements of Y(X) - S to these newly introduced vertices in a one-to-one manner.

Thus, we have obtained all the nonempty subsets of  $X_i$  on the edges of the resulting graph, say,  $T_3$ . Hence,  $T_3$  is properly set colorable.

We conclude the paper with a generalization of the notion of strong (proper) set colorings in the special case when the defining set X is taken to be a subset of the set N of nonnegative integers (or, for that matter, any linearly ordered set in place of N): An injective set assignment  $f : V(G) \cup E(G) \rightarrow 2^X$ ,  $X \subseteq N$  (or  $f: V(G) \rightarrow 2^X$ ) of a (p, q)-graph G is called a k-semi-strong (or k-semi-proper) set coloring of G if it satisfies the following conditions:



Fig. 4. Construction of a properly set colorable tree.

- (i)  $f^{\oplus}(uv) = f(u) \oplus f(v), \forall uv \in E(G),$
- (ii)  $f(G) \cup f^{\oplus}(G) = \{A_1, A_2, \dots, A_{p+q}\} \text{ (or } f^{\oplus}(G) = \{A_1, A_2, \dots, A_q\} \text{) where:}$ 
  - (a)  $A_1 < A_2 < \cdots < A_{p+q}$  with "<" defined on 2<sup>X</sup> by setting  $A < B \Leftrightarrow A, B \in 2^X$ , either |A| < |B| or
    - |A| = |B| and  $\min(A B) < \min(B A)$ ,
    - (b) for any  $A \in 2^X$ , if  $A_i < A$  and  $A < A_j$  for  $i < j, i, j \in \{1, 2, 3, ..., p+q\}$  then  $A = A_m$  for some  $m \in \{1, 2, 3, ..., p+q\}$ ,

(iii) 
$$|A_1| = k$$
.

The graph *G* is called a *k*-SSS graph if it admits a *k*-semi-strong set coloring and is called a *k*-SPS graph if it admits a *k*-semi-proper set coloring. In particular, a 1-SSS (1-SPS) coloring of *G* is simply called an SSS (SPS) coloring of *G*. Obviously then, an SSS coloring *f* of *G* is a strong (proper) set coloring of *G* if and only if  $X \subseteq N$  and  $f(G) \cup f^{\oplus}(G) = Y(X)$  ( $f^{\oplus}(G) = Y(X)$ ). Fig. 5 displays *k*-SSS (*k*-SPS) graphs for some values of *k*.

For any graph *G*,  $\beta_k(G)$  will denote the least cardinality of a set  $X \subseteq N$  with respect to which *G* has a *k*-SSS coloring. From the very definition, it follows that for any *k*-SSS coloring *f* with respect to a set



Fig. 5.

X of a graph G one must have

$$p+q \leq 2\beta_k - \sum_{j=0}^{k-1} {\beta_k \choose j} - L_k,$$

where  $L_k$  is the number of k-subsets of X which do not belong to  $f(G) \cup f^{\oplus}(G)$ . Furthermore, the bound is best possible.

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